

# Scientific Computing

## Announcements

→ Homework 4 assigned, see D2L  
Due Wednesday, April 2, 11:59pm

Today

→ Hill Climbing

March 24, 2025

Office Hours:

Mon + Fri

9:30am - 10:30am

Cudahy 307

## Problem Setup:

- \* Search space  $S$  full of candidates / possible solutions

- \* Scoring function: score( $x$ ),  $x \in S$   
(also called "fitness" w/ biological inspiration, or "quality")

\* A way to generate either

$x$  = current candidate  
= where you're standing in the mountains

- all the candidates "near" a candidate, the "neighborhood"  $\text{nbhd}(x)$

probably doesn't  
make total  
sense in  
cns space

OR

one

- a random candidate near a candidate (sometimes called a "tweak")  $\text{tweak}(x)$ .

"nearby" is up for you to define, and different definitions can totally change how a metaheuristic behaves.

Two running examples in this section.

1) TSP:

\* discrete — finite search space

\* score = cost of tour, want to minimize

\* nbhd(x): Suppose  $x = C_1 \rightarrow C_2 \rightarrow \dots \rightarrow C_n \rightarrow C_1$

"n-1 choose 2" Define the neighborhood to be all ways of picking two cities and swapping them (excluding  $C_1$ ).  $\approx n^2$

$$\text{size} = \binom{n-1}{2} = \frac{1}{2}(n-1)(n-2) \quad (\text{big!})$$

\* tweak(x): a random thing in the nbhd of x

(2) optimizing a continuous function  
in two variables  $f(x,y)$ .

\* continuous - infinite

\* score = value of the function

\* nbhd(x) = all points within some  
fixed distance  $\delta$  of x

small #

\* tweak(x) = a random point in nbhd(x)

infinite

# MH #1: Random Search

best = random element of  $S$

while True: (quit whenever you want)

$x$  = random element of  $S$

    if  $\text{score}(x) > \text{score}(\text{best})$ :

        best =  $x$

Possible stopping conditions:

- \* best score does not improve for  $N$  iters
- \* preset number of iters
- \* you get impatient

Gradient Ascent inspires this next one.

MH #2: Steepest Ascent Hill-Climbing (Discrete only)

$x$  = random element of  $S$

while True:

$N = \text{nbhd}(x)$

$s$  = element of  $N$  with the best score

if  $\text{score}(s) > \text{score}(x)$ :

$x = s$

else:

quit

Stopping conditions:

\* run out of time

\* no further improvement

What does this do? Climbs right up the hill you start on.

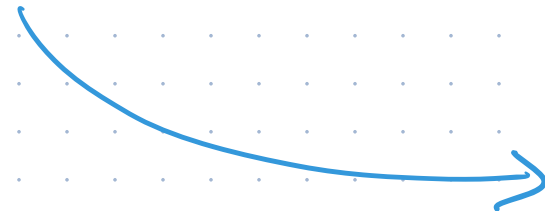
### Pros

- \* Finds a local optimum

### Cons

- \* Unlikely to find global optimum except in very nice spaces
- \* very slow, especially if nbhds are big, like TSP.

why?





Only really doing two things:  
(1) generating the neighborhood  
(2) scoring each element of it

Scoring a tour with 300 cities is not  
horrible - 300 distance calculations  
(two subtractions, two squarings,  
one addition, one square root)

But bad when you do it  
 $\binom{299}{2} = 44,551$  times.

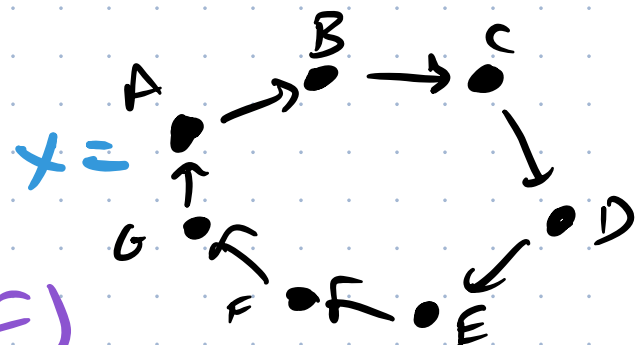
Often, you don't have to restore a solution from scratch because it only changes a little bit. More on this in a bit.

Demos:    03 - TSP    St. Asc.    50  
              04 - TSP    St. Asc.    300    slow!

How can we speed up scoring? Think about our tweak function. Suppose we have a tour:

Let d = distance function.

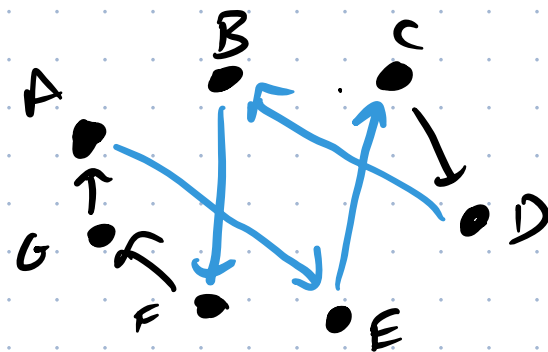
$$\text{Score} = d(A,B) + d(B,C) + d(C,D) + d(D,E) + d(E,F) + d(F,G) + d(G,A)$$



Swap B and E:

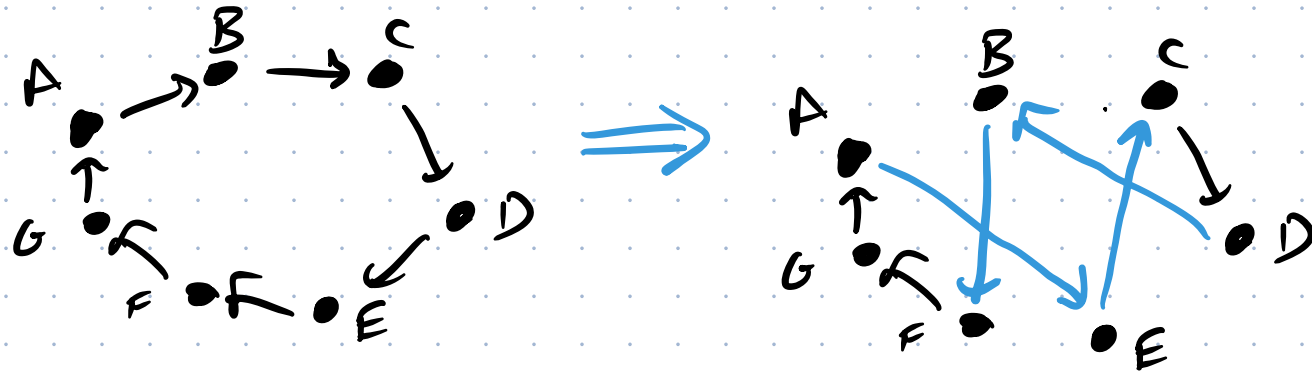
A → B → C → D → E → F → G → A

A → E → C → D → B → F → G → A



$A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow A$

$A \rightarrow E \rightarrow C \rightarrow D \rightarrow B \rightarrow F \rightarrow G \rightarrow A$



Four edges change.

$$\text{Score} = \cancel{d(A, B)} + \cancel{d(B, C)} + d(C, D) + \cancel{d(D, E)} + \cancel{d(E, F)} + d(F, G) + d(G, A) + \cancel{d(B, F)}$$

$d(A, E)$      $d(E, C)$      $d(C, D)$      $d(D, B)$

If you have 300 cities, still only 4 edges change:

new score = old score - 4 edges + 4 edges.

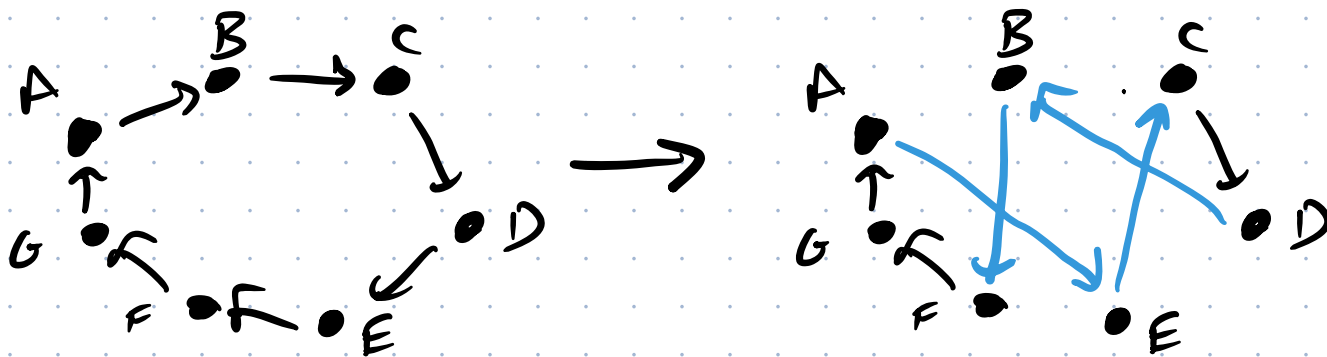
8 distances instead of 300 =  $\frac{300}{8} = 37.5\times$  faster!

Demo 05

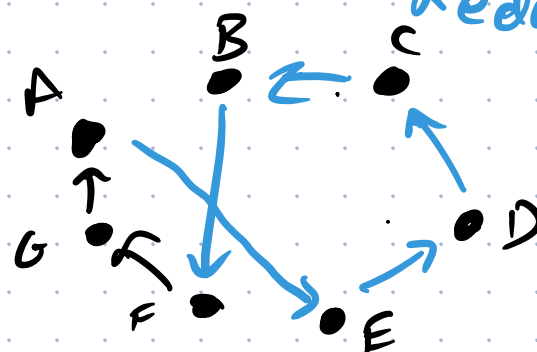
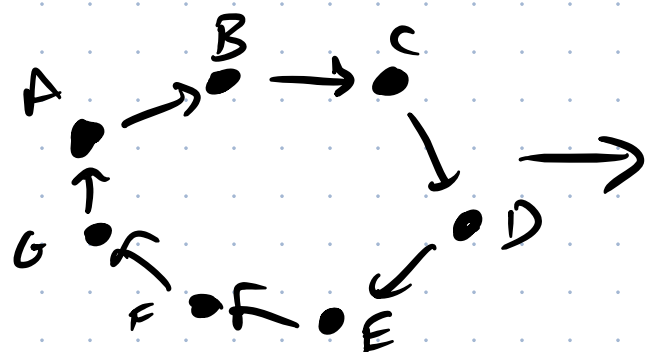
Second question: Is this a good tweak?

Our demos suggest maybe not.

Why not? Small tweaks are better (at least for now).



4 edges changed... can we change just 1?  
No. How about 2?



4 edges: swap 2 cities  
2 edges: reverse block

Assuming distance is symmetric yes.

This is picking two cities and reversing the whole block.

A → B → C → D → E → F → G → A

A → E → D → C → B → F → G → A  
(faster scoring too!)

Demos: 06 - SA RB 50  
07 - SA RB 300

08 / Fast Score

The big theme of MTHs is that they are super flexible. You can and should tweak them in all kinds of ways for particular problems. Always experiment.



How can we adapt this for continuous spaces?

How can we adapt this for continuous spaces?

Check  $n$  things in the neighborhood and take the best.

# MH #3 n-Trial Steepest Ascent

$x$  = random element of  $S$

while True:

temp =  $x$

repeat  $n$  times:

$s$  = tweak( $x$ )

if  $\text{score}(s) > \text{score}(\text{temp})$

temp =  $s$

$x$  = temp

(if nothing  
beats  $x$ , it  
stays the same)

temp is the best neighbor out  
of the  $n$  we tried

Later we will see good ways to tweak for  
continuous spaces.

When  $n=1$ , this is just called "Hill Climbing"

MH #4: Hill Climbing

$x$  = random element of  $S$

while True:

$s = \text{tweak}(x)$

if  $\text{score}(s) > \text{score}(x)$ :

$x = s$

Repeat: small

→ Take a step.

If better, stay.

If worse, go back.

Repeat

Demos: 09 - TSP HC Swap 2 50  
 10 - TSP HC Swap 2 300  
 11 - TSP HC RB 50  
 12 - TSP HC RB 300

HC = Hill Climbing  
 SA = Steepest Ascent

RB = reverse a whole  
 block of cities

50 cities

SA Swap 2 9.878

SA RB 6.487

HC Swap 2 8.428

HC RB 6.412

300

cities

SA Swap 2 32.828 cities

SA RB 14.362

HC Swap 2 29.439

HC RB 14.252

Regular HC beat Steepest Ascent  
 RB very much beat Swap 2