Wed, Apr 21 HW 6 assign. Fri, due last day of class Optional assign. on last day, due during exam week HW 5: blue lights graph part. 13005 good sol: 215 best sol: 21.2514 1351 21.25139 blue lights: mult: ple tweaks prick a random light replace it with a new random light a lot of exploration not great for exploitation tweaks = picking a random solution GP unpenalized score

allow the two sides to have different sizes, but penalize it $\left[\frac{502 - 498}{5000}\right] = 4$ $5000 \cdot (1 + \frac{4}{1000})$ Topic 18 - Variations on Local Search "Local Search" = MHs where we look nearby, like H-C, S.A., T.S, etc. Today: two variations to give you a taste of what's possible #1 Variable Neighborhood Search <u>Idea</u>: Define different kinds of tweaks (different neighborhoods). N,(x), Nz(x), Nz(x),, Nd(x) more and more dramatic Often, but not always: $N_1(x) \subseteq N_2(x) \subseteq N_3(x) \subseteq \dots \subseteq N_d(x)$



With the blue lights: N, = more one light a tiny amount Nz = delete a light and randomly. re place it $N_{1}(x) \leq N_{2}(x)$

Pseudocode: Let x be a random solution. While True: Pick & randomly from N.(x) if s is better than x: keep it and start the while loop over else: Pick & rondomly from Nz(x) if s is better than x: keep and stort over

else: Pick 5 randomly from Nz(x) and so on If you go through all neighborhoods with no improvement, just start over with $N_{1}(\varphi)$ * You can try a neighborhood more than once before moving onto the next Gne. We need: * a neighborhood structure Ex: Knopsack NK(X) = delete k random items, then greedily add items back until full Ex: "Metaheuristras: From Design to Implementation" by Talbi

Example 2.35 Many neighborhoods in a real-life application. This example illustrates the use of many neighborhoods (k = 9) to solve a scheduling problem [26]. Many oil wells in onshore fields rely on artificial lift methods. Maintenance services such as cleaning, stimulation, and others are essential to these wells. They are performed by workover rigs. Workover rigs are available in a limited number with respect to the number of wells demanding service. The decision which workover rig should be sent to perform some maintenance services is based on factors such as the well production, the current location of the workover rig in relation to the demanding well, and the type of service to be performed. The problem of scheduling workover rigs consists in finding the best schedule S^* for the available *m* workover rigs, so as to minimize the production loss associated with the wells awaiting service. A schedule is represented by an ordered set of wells serviced by workover rigs. A variable neighborhood search metaheuristic has been proposed for this problem using the following neighborhoods in the shaking procedure:

- 1. Swapping of routes (SS) where the wells associated with two workover rigs are exchanged.
- 2. Swapping of wells from the same workover rig (SWSW) where two wells serviced by the same workover rig are exchanged.
- 3. Swapping of wells from different workover rigs (SWDW) where two wells affected by two different workover rigs are exchanged.
- 4. Add/drop (AD) where a well affected by a workover rig is reassigned to any position of the schedule of another workover rig.
- 5. Two applications of the SWSW transformation.
- 6. Two applications of the SWDW transformation.
- 7. Three applications of the SWDW transformation.
- 8. Successive application of two AD transformations.
- 9. Successive application of three AD transformations.

Ex 3: TSP Good tweak we've used: "pick two cities, and reverse the path in between". pictuse:

Another way to phrase this: "delete two random edges, then reconnect the tour in the cheapest way that is still a valid cycle not a valid TSP solution "k-opt": delete k edges, consider all Valid ways of reconnecting

the tour, pick the cheapest 3-opt: most popular for H-C k=3: $(\mathbf{x}) (\mathbf{x}) (\mathbf{y}) (\mathbf{y}) (\mathbf{y}) (\mathbf{y})$ 8 volid possibilities to check.